



Exercice 5.3

- $p \wedge q$
- $p \vee (q \wedge r)$
- $q \wedge r$
- q, r, p
- $\neg(p \vee (q \wedge r)) \Rightarrow (p \wedge q)$

2. $Sf(P) = \{P\}$ si P atomique

$$Sf(\neg P) = \{\neg P\} \cup Sf(P)$$

$$Sf(P \circ Q) = Sf(P) \cup Sf(Q) \cup \{P \circ Q\} \quad \text{où } \circ \in \{\wedge, \vee, \Rightarrow\}$$

3.]

$\phi(P)$: si P a n connecteurs alors $|Sf(P)| \leq 2n+1$

$$\rightarrow Sf(P) = \{P\} \quad \text{où } P \text{ atomique}$$

$$|Sf(P)| \leq 2 \cdot 0 + 1 = 1 \quad \text{par hyp. de rec.} \quad \checkmark$$

$$\rightarrow Sf(\neg P) = Sf(P) \cup \{\neg P\}$$

$$|Sf(\neg P)| = |Sf(P)| + 1 \leq 2n+1+1 = 2(n+1) \leq 2(n+1)+1$$

si P a n connecteurs, donc $\neg P$ a $n+1$ connecteurs

$$\rightarrow Sf(P \circ Q) = Sf(P) \cup Sf(Q) \cup \{P \circ Q\}$$

$$\text{donc } |Sf(P \circ Q)| = (2|Q|+1) + (2|R|+1) + 1 \quad \text{par HR}$$

$$\leq 2(|Q|+|R|) + 3$$

$$\leq 2(|Q \circ R|) + 1$$

atteint lorsque les connecteurs ont l'arité 2.

Exercice 6.1

$$1. \forall x_1, \dots, x_{n+1}, \bigvee_{1 \leq i < j \leq n+1} x_i = x_j$$

est exactement l'axiome des modèles équationnels avec moins de n éléments.

ou bien

$$\exists x_1, \dots, x_n, \left(\bigwedge_{i < j} x_i \neq x_j \wedge \left(\bigvee_{1 \leq i \leq n} x_i = z \right) \right)$$

Exercice 6.2

1) Soient $x_1, \dots, x_n \neq 0$ $x_n = S(x_{n-1}), \dots, x_2 = S(x_1)$

$$\{ S^k(0_N) \in \mathbb{N} \mid k \in \mathbb{N} \} \subseteq \mathbb{N}$$

Si le domaine serait fini, on aurait :

$$S_N^p(0_N) = S_N^{n+p}(0_N)$$

$$\text{donc } 0_N = S^p(0_N) = S(\overbrace{S^{p-1}(0_N)}^{\neq 0}) = 1_{0_N}$$

contradiction

injectivité + non-surjectivité $\Rightarrow$

2) Oui.

3) Non, on peut avoir $0 = S(0) = 1$
 $1 = S(1)$

et le domaine est $\{0, 1\}$

exercice 6-5

I

$\mathcal{T}(1) = (\text{faux}, \text{faux})$
 $\mathcal{T}(2) = (\text{vrai}, \text{faux})$
 $\mathcal{T}(3) = (\text{faux}, \text{vrai})$
 $\mathcal{T}(4) = (\text{vrai}, \text{faux})$
 $\mathcal{T}(5) = (\text{faux}, \text{faux})$
 $\mathcal{T}(6) = (\text{vrai}, \text{vrai})$

25 4 valeurs différentes

$$\begin{aligned} \min &= 1 \\ \max &= 9 \end{aligned}$$

3) $\phi(A) \stackrel{\text{def}}{=} \mu + i, i'$ si $i \cong i'$ alors $\text{val}(i, A) = \text{val}(i', A)$

$$\forall x, A \quad i, i' \quad i \simeq i' \quad \text{val}(i, \forall x, A) = \text{val}(i', \forall x, A)$$

$$L \in \mathcal{D} \quad \text{val}(i + (x \mapsto d), A) = \text{val}(i' + (x \mapsto d), A)$$

$$i + (x \mapsto d) \simeq i' + (x \mapsto d)$$

$$\begin{array}{l} A \cap B \\ 7A \end{array} \quad \begin{array}{l} \therefore i' \quad i \simeq i' \quad \text{val}(i, 7A) = \text{val}(i', 7A) \\ \text{per KR} \quad \text{val}(i, A) = \text{val}(i', A) \end{array}$$

$$P(x) \quad i \approx i' \quad \text{val}(i, P(x)) = \text{val}(i', P(x))$$

$$Q(x) \quad i \approx i'(x) \quad \text{vrai par def} \quad \text{le } \approx$$

4) (a) image invers, l'ensemble fini

(f) $\bar{1}, \bar{2}, \bar{3}, \bar{6}$

(c) triviale

Exercise 6.6

1.) $\neg P \Leftrightarrow \exists F(P, \perp, \top)$

$$P \wedge Q \Leftrightarrow \exists F(P, Q, \perp)$$

$$P \vee Q \Leftrightarrow \exists F(P, \top, Q)$$

$$\neg P \vee Q \Leftrightarrow \exists F(\neg P, \top, Q)$$

2.)

3.)

4.) $\text{norm}(X) = X$
 $\text{norm}(T) = T$
 $\text{norm}(\perp) = \perp$
 $\text{norm}(\exists F(A, B, C)) = \exists F_n(A, \text{norm}(B), \text{norm}(C))$

$$\exists F_n(X, Q, R) = \exists F(X, Q, R)$$

$$\exists F_n(T, Q, R) = \exists F(T, Q, R)$$

$$\exists F_n(\perp, Q, R) = \exists F(\perp, Q, R)$$

$$\exists F_n(\exists F(A, B, C), Q, R) = \exists F_n(A, \exists F_n(B, Q, R), \exists F_n(C, Q, R))$$

Exercise 6.7

